

# CS 350S: Privacy-Preserving Systems

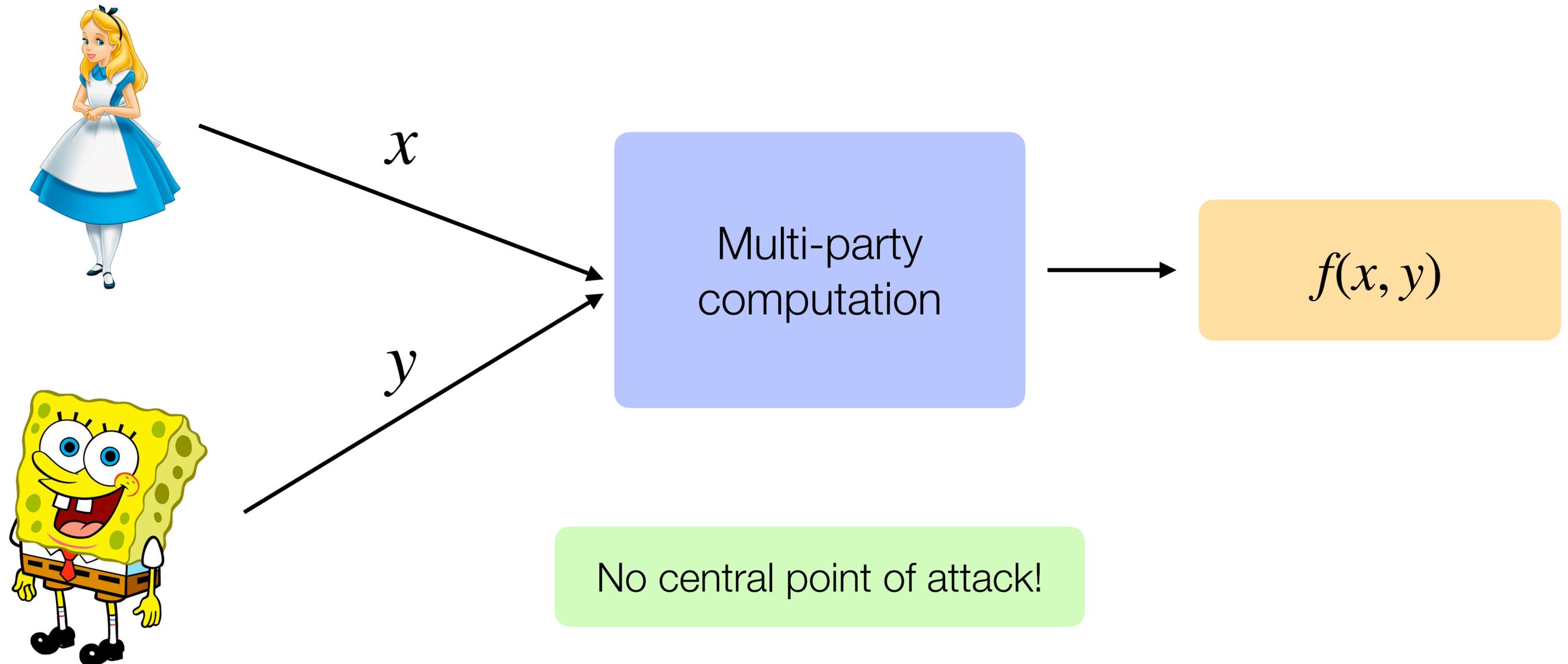
Differential privacy

# Outline

1. Differential privacy definition
2. Differential privacy mechanism
3. Differential privacy applications
4. Logistics
5. Student presentation

# MPC recap

(Informal) Any computation that can be performed with a trusted third party can be securely computed *without* one!



# Use case: private aggregate statistics

Aggregation function  $f$

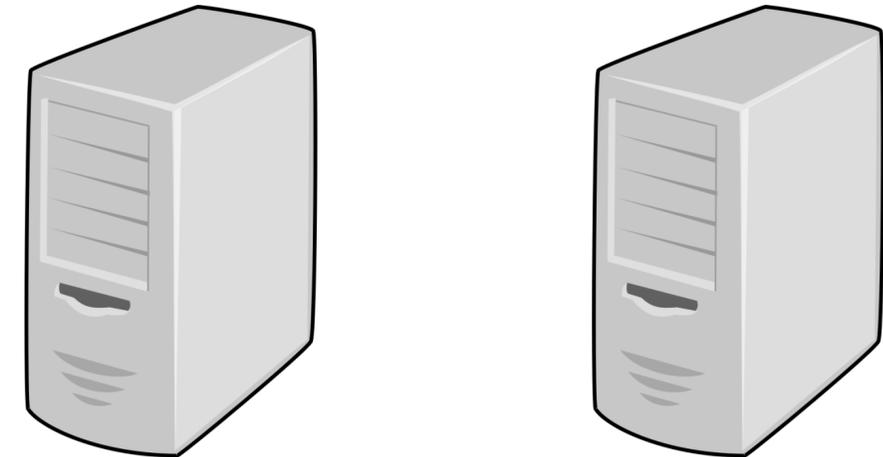
**Correctness:** If all servers are honest, servers learn  $f(x_1, x_2, \dots, x_n)$

**Privacy:** If one server is honest, servers only learn  $f(x_1, x_2, \dots, x_n)$

- Privacy with 1 malicious server

**Robustness:** Malicious clients have bounded influence

$$f(x_1, x_2, \dots, x_n)$$



Examples of when the output of a computation can reveal private information?

# Anonymization is not enough (AOL query dataset)

- AOL query dataset had >20M anonymized search queries from 650,000 AOL users over 3 months
- Dataset released where each username was replaced with a random identifier
- Queries for
  - “Landscapers in Lilburn, Ga”
  - Several people with last name Arnold
  - “Homes sold in shadow lake subdivision Gwinnett county Georgia”
  - ... other sensitive queries
- Only 14 citizens with last name Arnold in Gwinnett County
- Found that user was Thelma Arnold, 62-year old woman in Georgia



# Anonymization is not enough (Netflix prize dataset)

[Narayanan, Shamtikov]

- Netflix prize dataset has anonymized movie ratings for 50K Netflix subscribers
- Adversary who has a small amount of background knowledge (e.g., dates of some ratings within a 14-day window, some approximate / potentially incorrect ratings), can uniquely identify a record in the published dataset
- Used IMDB ratings that users posted publicly under their own name
- A seemingly “harmless” dataset without identifying information can be a privacy violation

**Robust De-anonymization of Large Datasets  
(How to Break Anonymity of the Netflix Prize Dataset)**

Arvind Narayanan and Vitaly Shmatikov

The University of Texas at Austin

# Aggregating alone is not enough

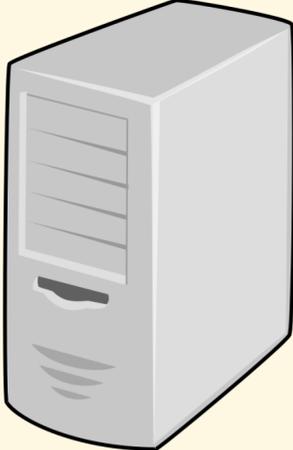
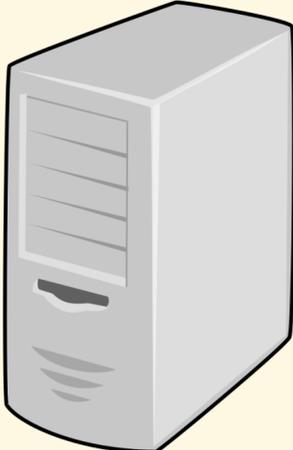
Number of users with medical condition X

Day 1

Day 2



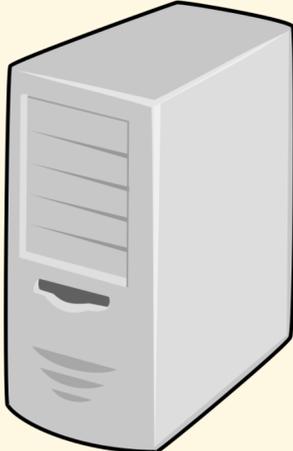
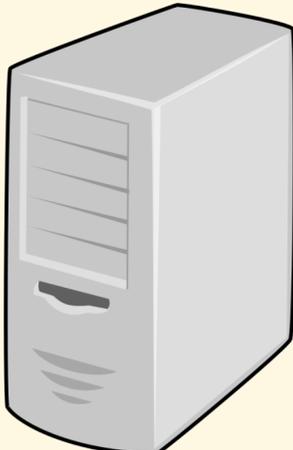
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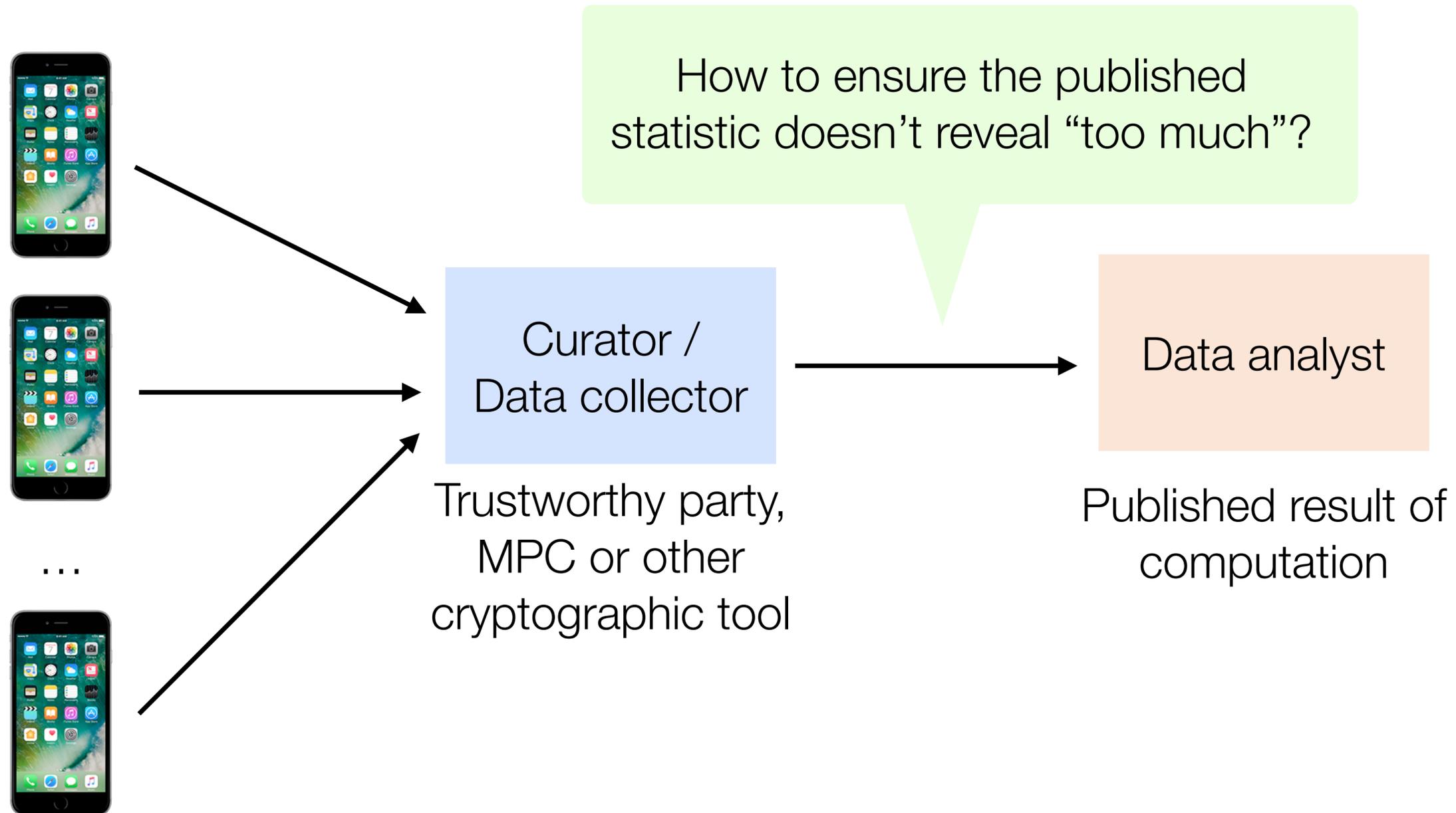


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# Building a private system

- Cryptographic protocols like MPC often focus on *how* to provide privacy
  - Given a function  $f$ , how to compute the output of the function while hiding the inputs and intermediate data
- Today: What function  $f$  should we be computing to ensure privacy?
  - How to define privacy?
  - How to provide privacy?
- Systems need both

# Model



# Attempt #1 at defining privacy

The result of the analysis should not allow an adversary to learn any information about individual users' data that it would not learn without access to the result.

Does this work?

No — the revealed statistic needs to be useful, and so we cannot use some sort of cryptographic definition where the output appears random and reveals nothing about the inputs

# Attempt #2 at defining privacy

The result of the analysis should not allow an adversary to learn any information about any individual in the dataset that the adversary could not learn if the individual was not in the dataset.

Does this work?

No — the revealed statistic needs to be useful, and so cannot be independent of all participants

# Attempt #2 at defining privacy

The result of the analysis should not allow an adversary to learn **any information** about **any individual** in the dataset that the adversary could not learn if the individual was not in the dataset.

Does this work?

No — the revealed statistic needs to be useful, and so cannot be independent of all participants

Can we weaken part of this statement?

# Attempt #3 at defining privacy

The result of the analysis should not allow an adversary to learn **any information** about **most individuals** in the dataset that the adversary could not learn if the individual was not in the dataset.

Does this work?

Kind of — possible to sample a small group of users and then compute statistics from only their data

... but some users have to give up their privacy

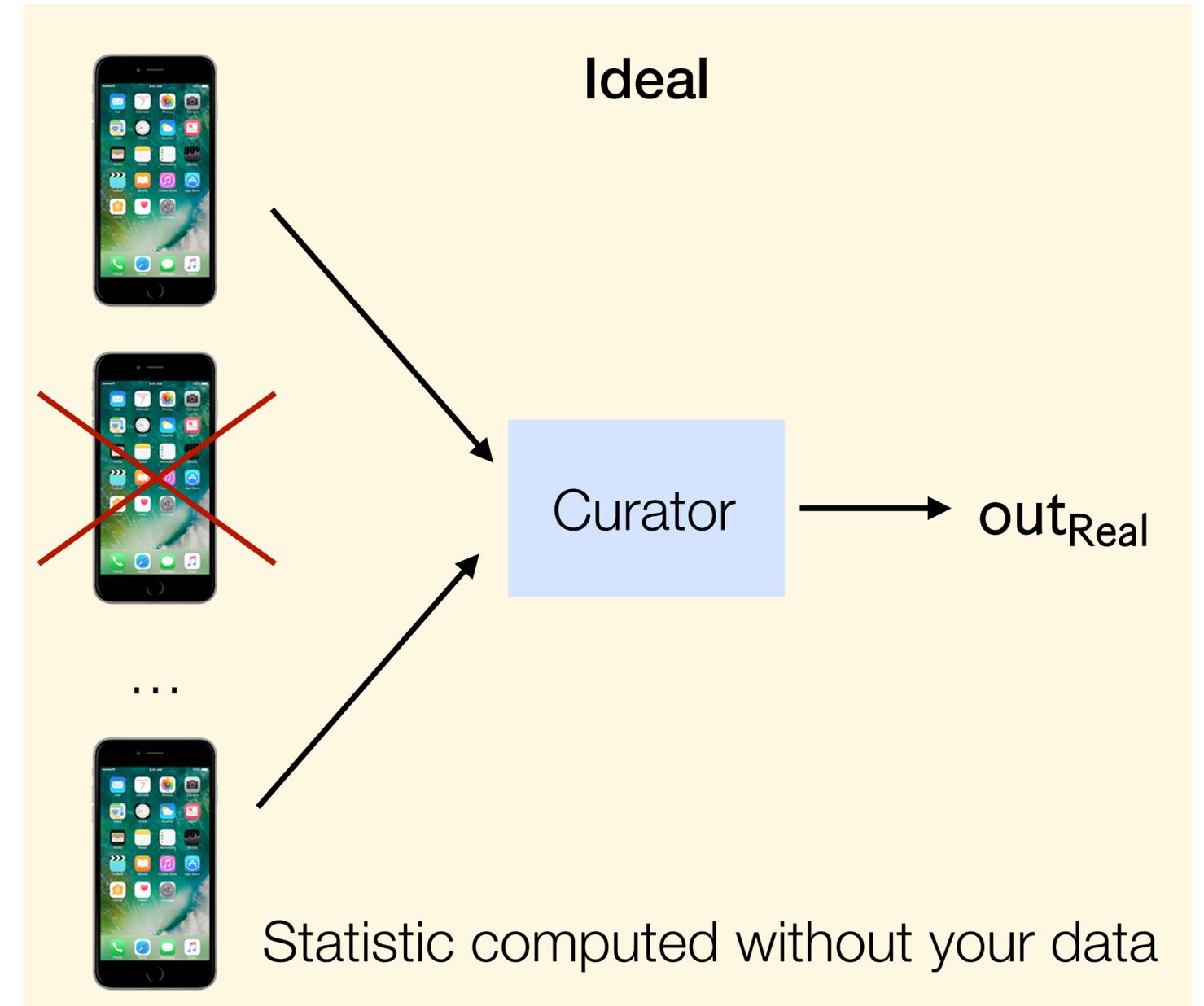
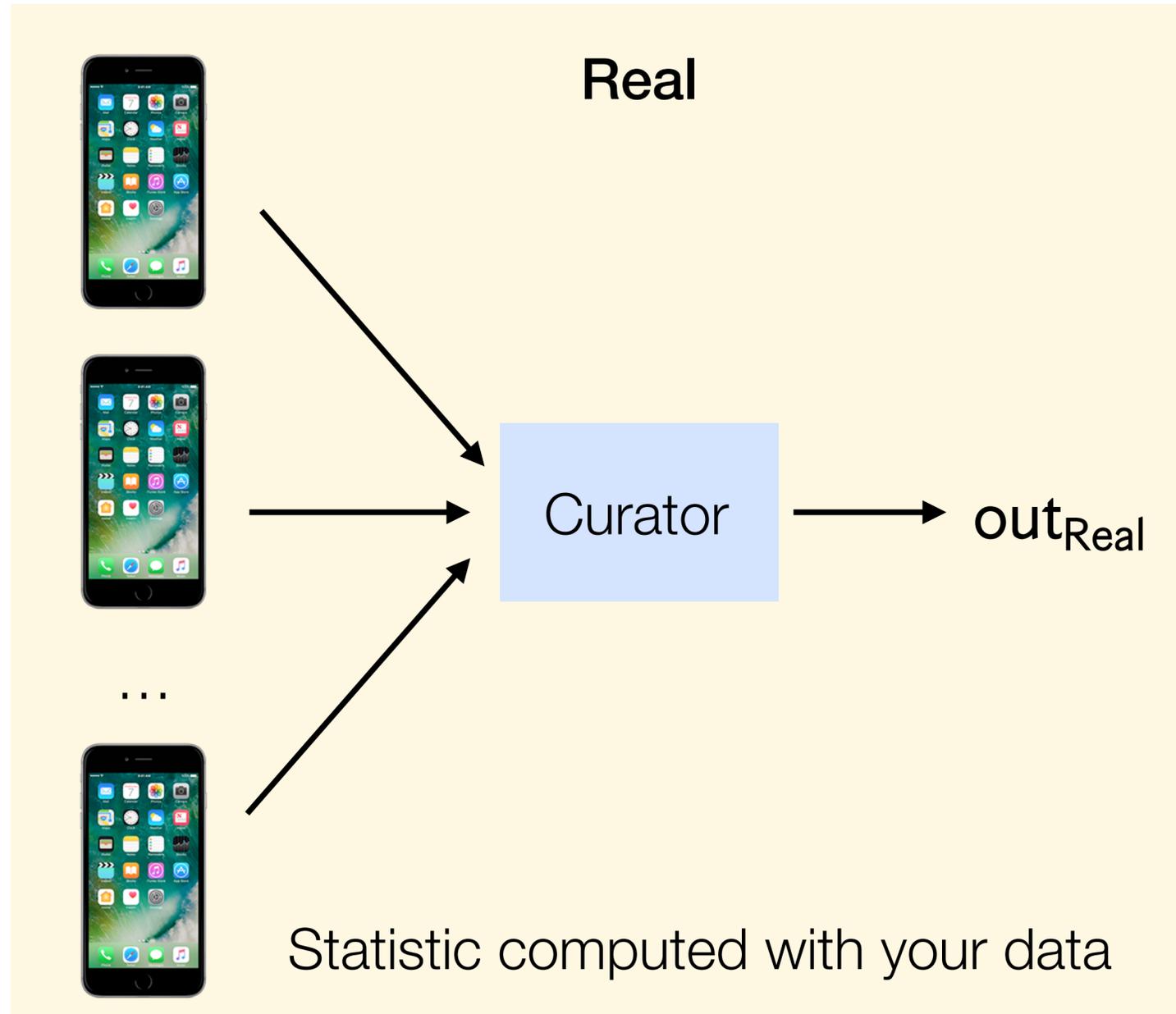
# Attempt #4 at defining privacy

The result of the analysis should not allow an adversary to learn **anything new confidently** about **any individual** in the dataset that the adversary could not learn if the individual was not in the dataset.

Does this work?

Yes, if we define what “anything new confidently” means — intuition behind differential privacy

# Differential privacy



$$out_{Real} \approx out_{Ideal} \quad (\text{Not cryptographic indistinguishability})$$

# Formalizing differential privacy

[Dwork, Sherry, Nissim, Smith]

Mechanism  $\mathcal{M} : \mathcal{X}^n \rightarrow \mathcal{Y}$

For database with  $n$  rows of type  $\mathcal{X}$  and output statistic  $\mathcal{Y}$

Two databases are “neighboring” if they differ in at most one row

A mechanism  $\mathcal{M}$  is  $\varepsilon$  differentially private if for all pairs of “neighboring databases”  $D, D'$  and every set of values  $S \in \mathcal{Y}$ :

$$\Pr[\mathcal{M}(D) \in S] \leq e^\varepsilon \cdot \Pr[\mathcal{M}(D') \in S]$$

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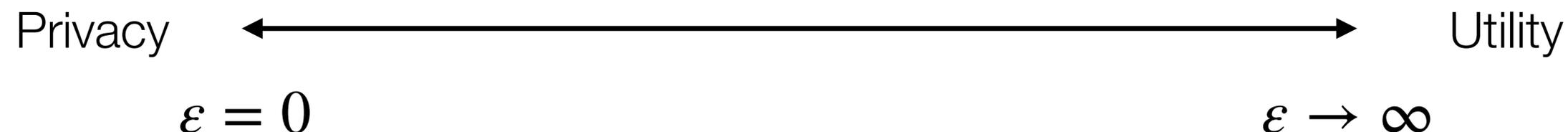
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How to set  $\epsilon$ ?



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Intuition: Any “bad event” when Alice is in the database would have happened with similar probability if she was not in the database

# Formalizing differential privacy

[Dwork, Sherry, Nissim, Smith]

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*Question: What are some databases where, even if we apply differential privacy, the published statistic still reveals some sensitive information about a user?*

# Example

Phone number	Timestamp
123-456-7890	11/17 8:00AM
911	11/17 9:00AM
123-456-7890	11/17 9:30AM
...	...

Will a  $\epsilon$  differentially private mechanism over this database hide the following with  $\epsilon$ -DP:

- Frequent calls to Bob (123-456-7890)? **X**
- One call to 911? **✓**

# Properties of differential privacy

**Post-processing:** Let  $\mathcal{M} : \mathcal{X}^n \rightarrow \mathcal{Y}$  be  $\epsilon$  differentially privacy and let  $f : \mathcal{Y} \rightarrow \mathcal{Z}$  be any (randomized) function. Then  $(f \cdot \mathcal{M}) : \mathcal{X}^n \rightarrow \mathcal{Z}$  is  $\epsilon$  differentially privacy.

- Intuition: Whatever computation is performed on the results of a differentially private query, the answer is still differentially private
- Implies resilience to any side information

**Composition:** Let  $\mathcal{M}_1, \mathcal{M}_2, \dots, \mathcal{M}_n$  be mechanisms where  $\mathcal{M}_i : \mathcal{X}^n \rightarrow \mathcal{Y}_i$  is  $\epsilon_i$  differentially private. Then  $\mathcal{M}(D) \rightarrow (\mathcal{M}_1(D), \mathcal{M}_2(D), \dots, \mathcal{M}_n(D))$  is  $\epsilon$  differentially

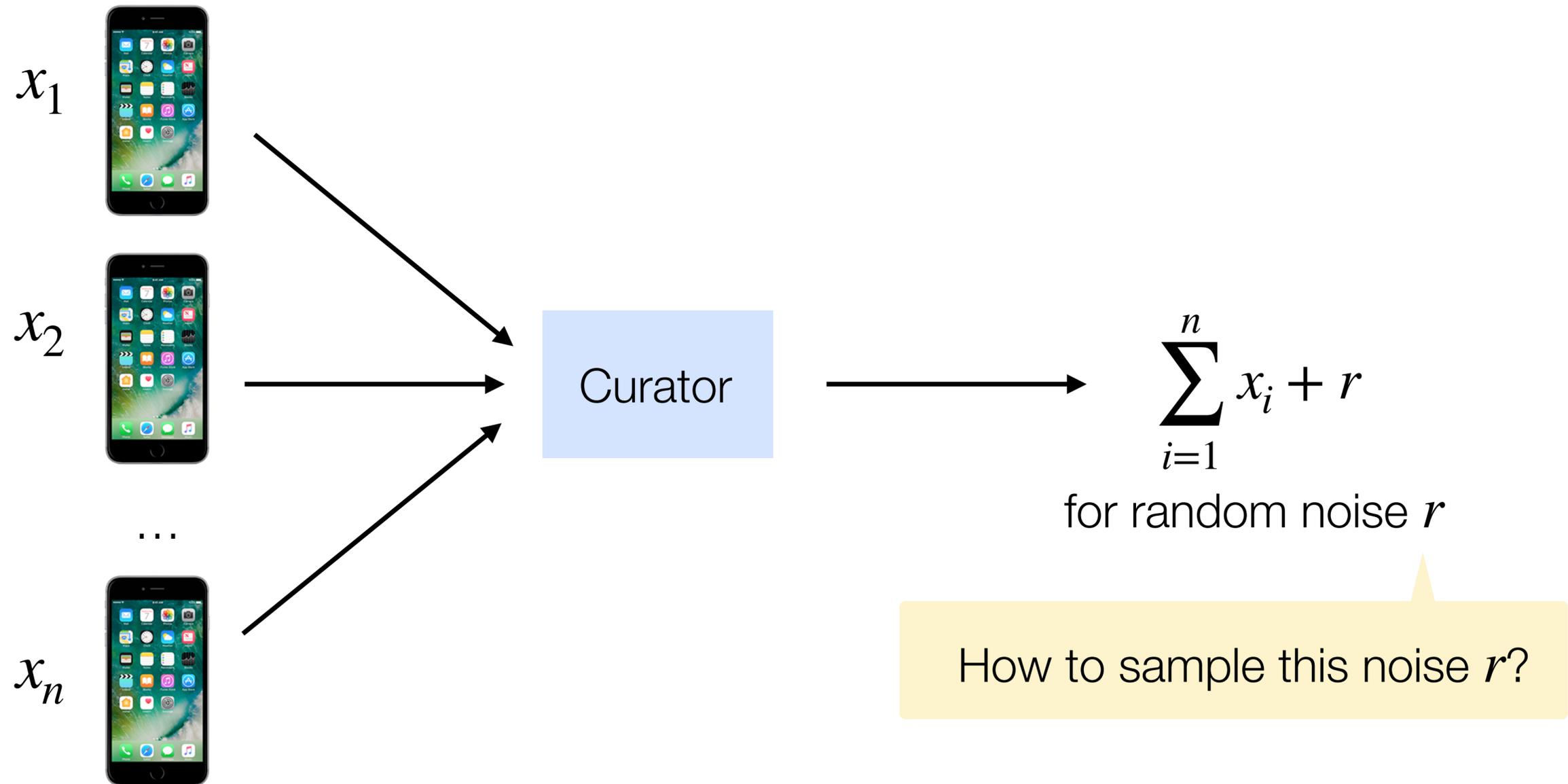
private for  $\epsilon = \sum_{i=1}^n \epsilon_i$

- Intuition: Answering more queries reveals more information
- Bounds overall privacy risk

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# Differential privacy at a high level



# Defining query sensitivity

**Sensitivity:** For a query  $q : \mathcal{X}^n \rightarrow \mathbb{R}$ , the sensitivity of  $q$  is  $\Delta q = \max_{D \sim D'} |q(D) - q(D')|$

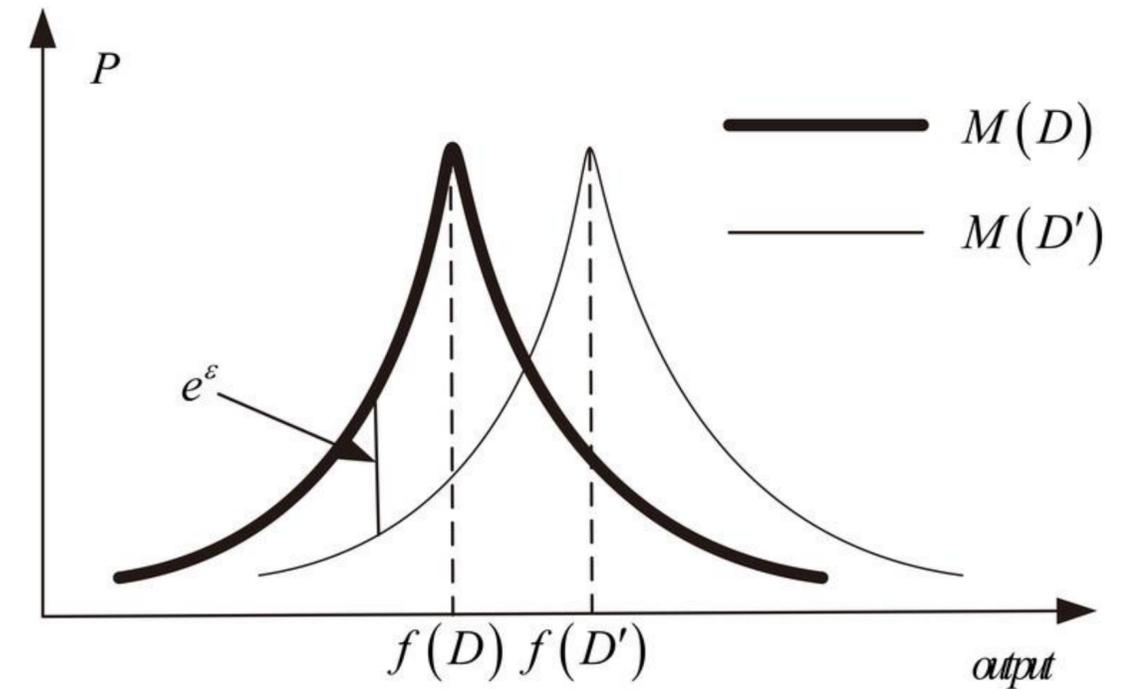
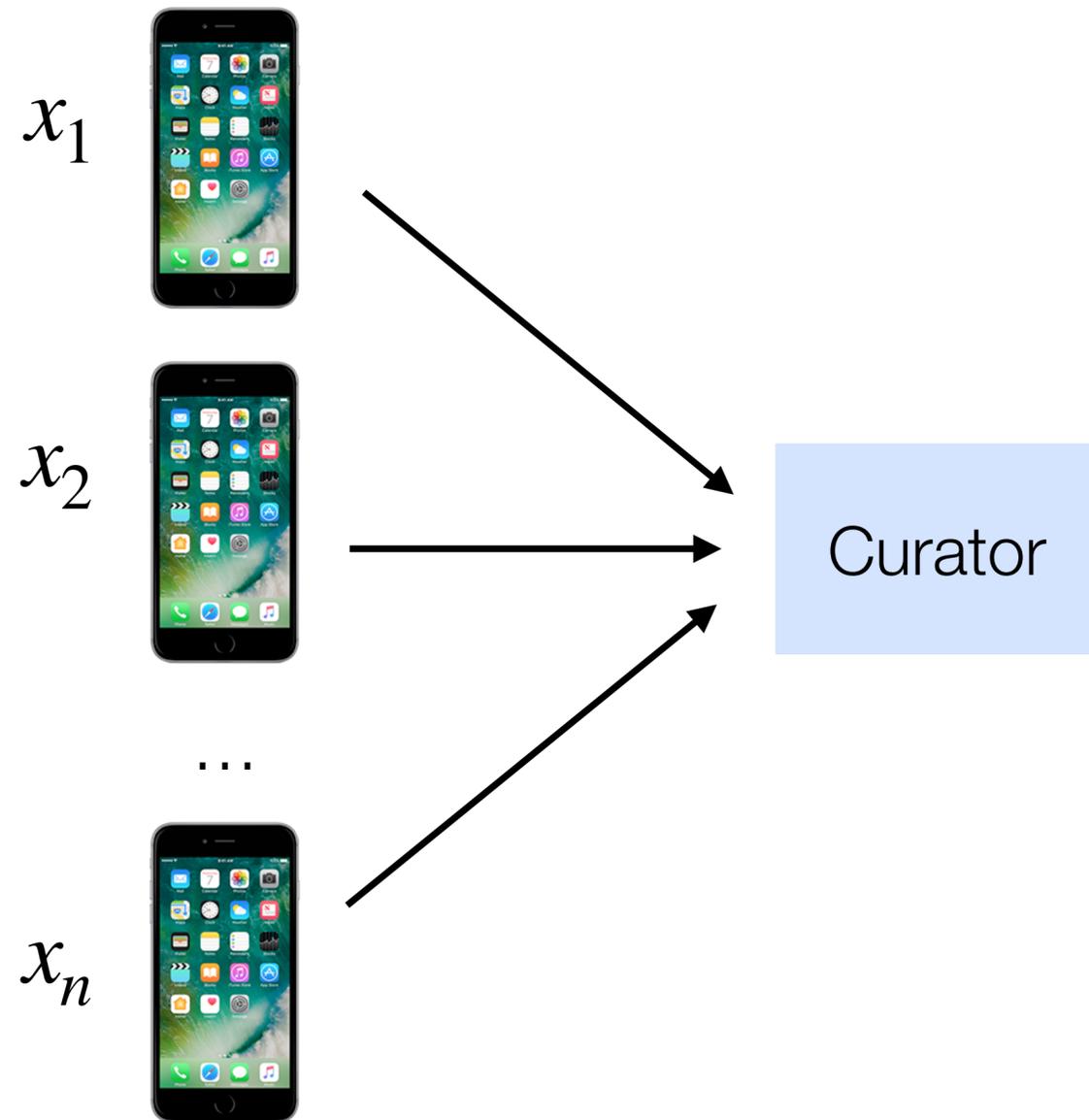
For counting query?

Sensitivity: 1

For sums of 8-bit values?

Sensitivity:  $2^8 = 256$

# Laplace mechanism



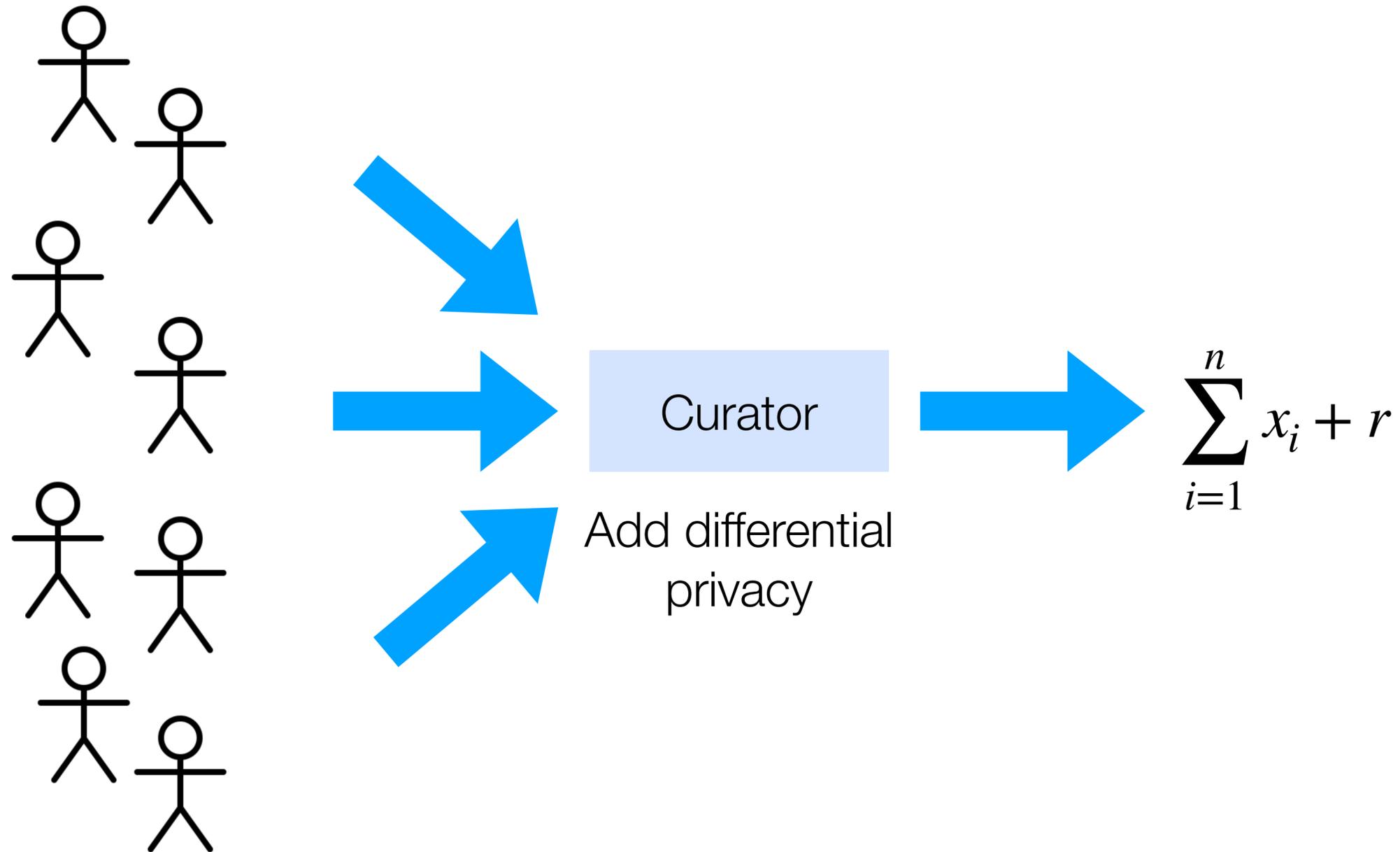
$$q(x_1, x_2, \dots, x_n) + r$$

Where  $r \sim \text{Lap} \left( \frac{\Delta q}{\epsilon} \right)$

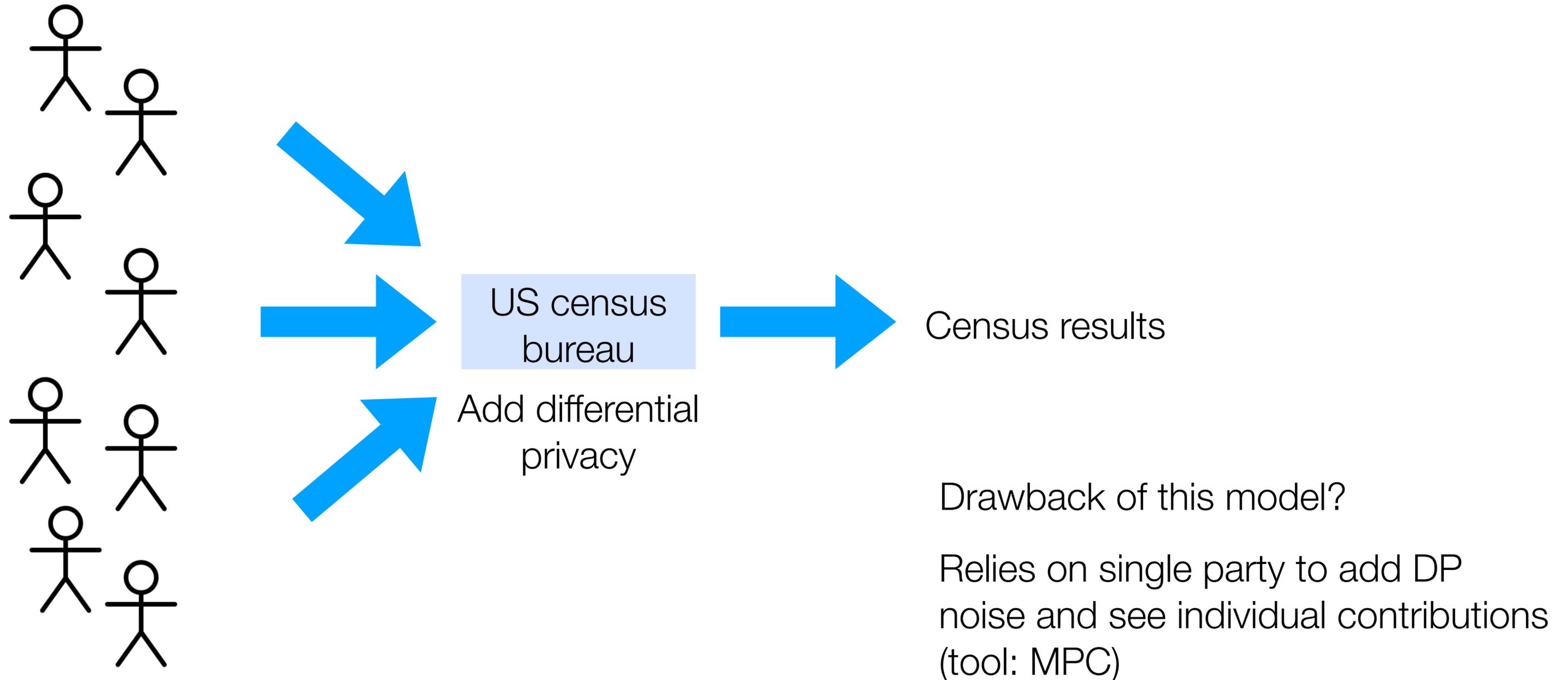
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# Central differential privacy

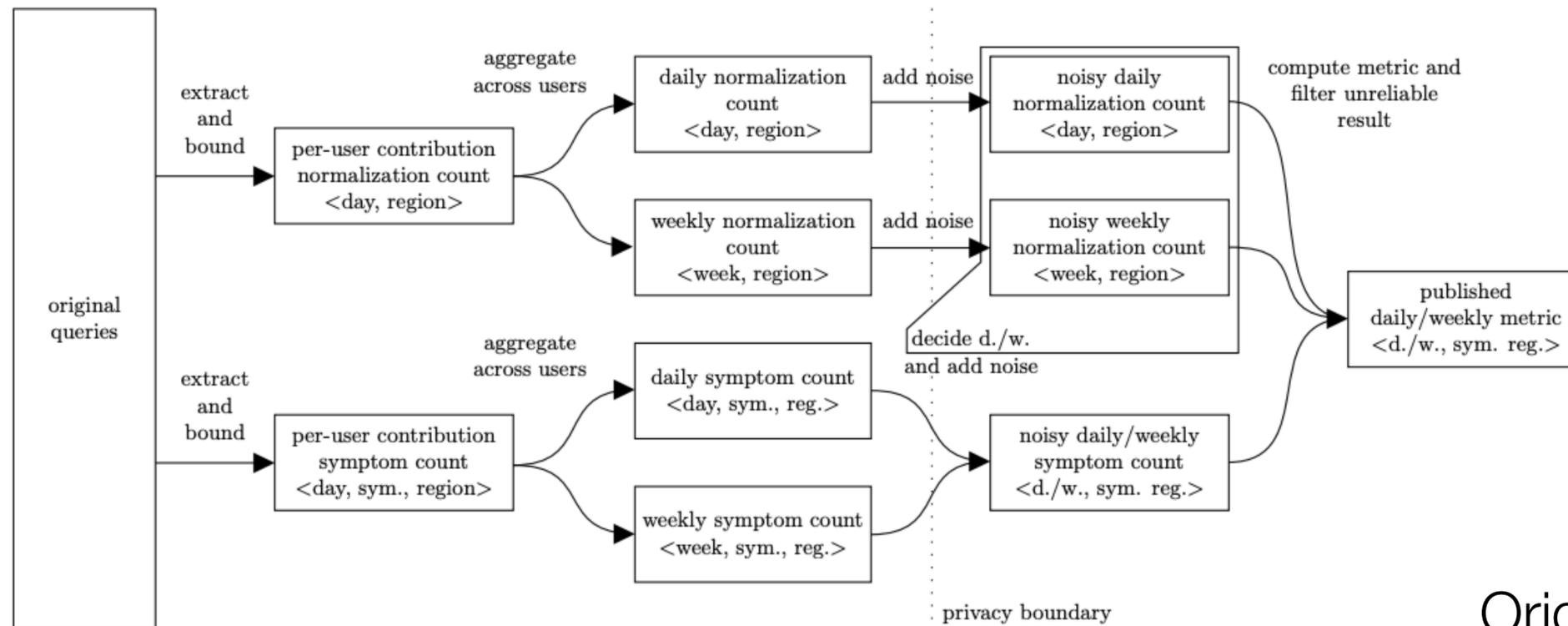


# Central differential privacy: US census



# Google COVID-19 search trends

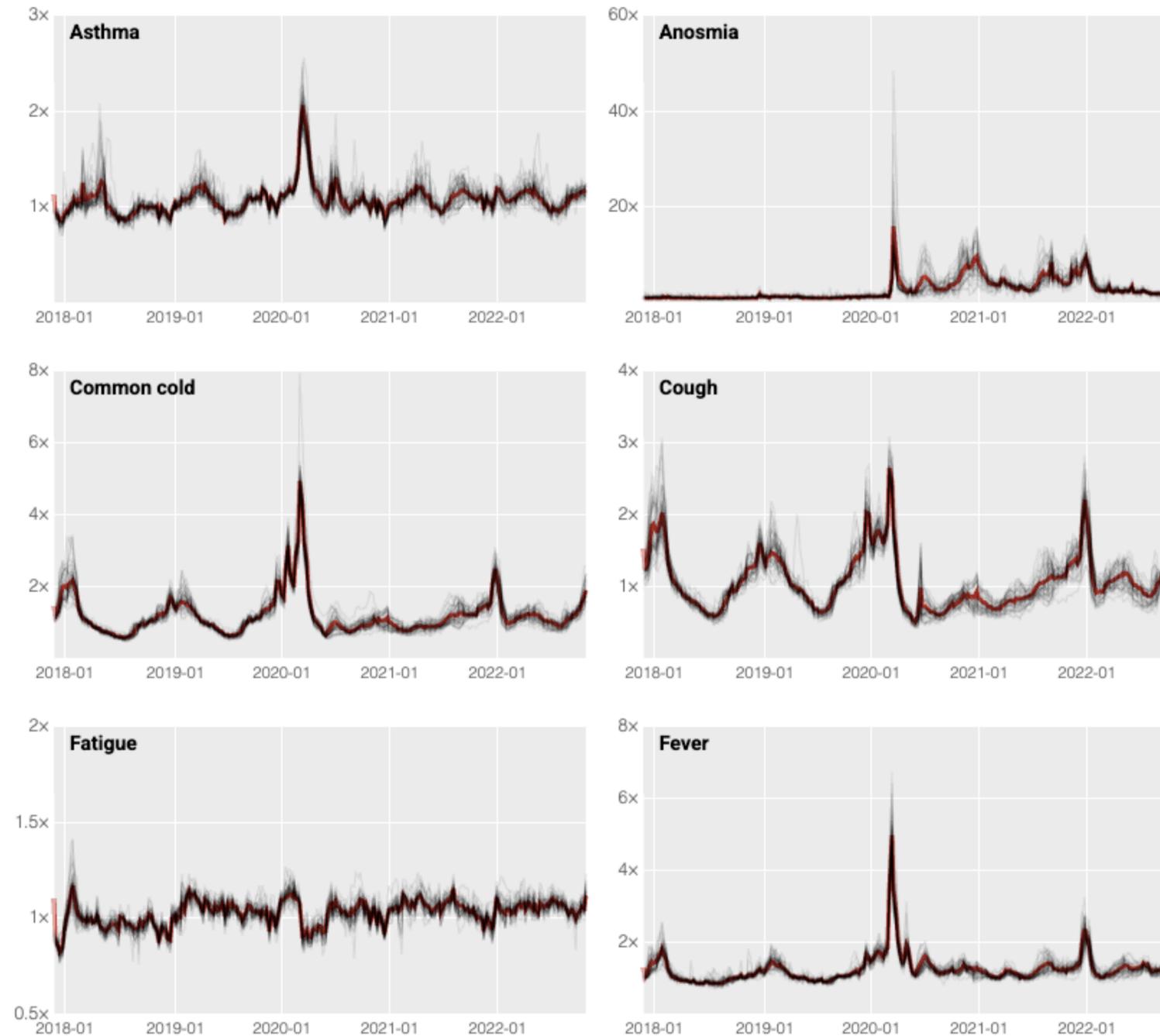
Also in the central DP model ( $\epsilon = 1.68$ )



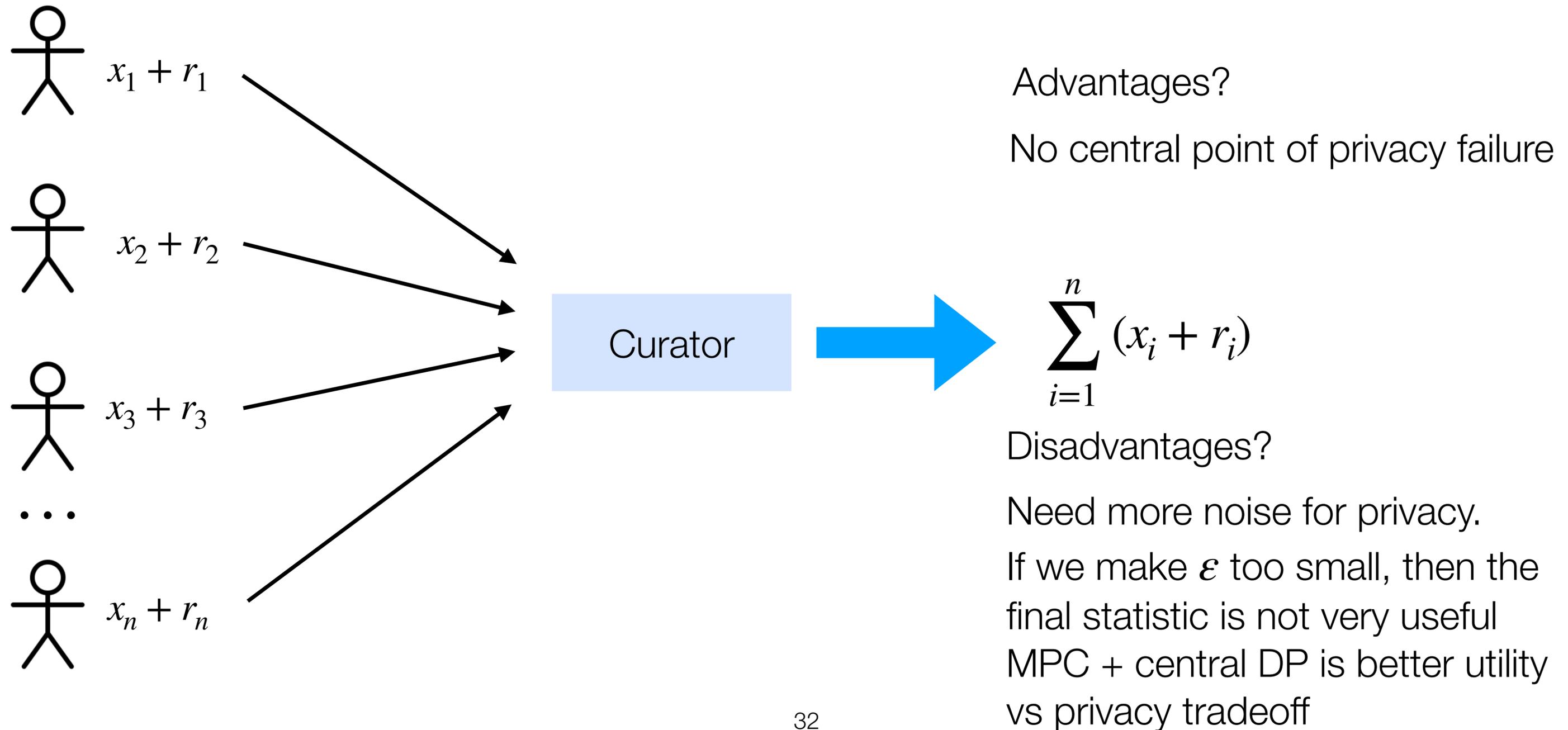
level	count	contribution	bound type
0	<2020-06-03, fever, United States>	1 (originally 4)	per-symptom
0	<2020-06-03, cough, United States>	1	
1	<2020-06-03, fever, California>	1 (originally 3)	per-symptom
1	<2020-06-03, fever, Nevada>	1	
1	<2020-06-03, cough, Nevada>	1	
2	<2020-06-03, fever, Santa Clara>	1 (originally 2)	per-symptom
2	<2020-06-03, fever, San Bernardino>	1	
2	<2020-06-03, fever, San Bernardino>	1	
2	<2020-06-03, fever, Clark>	1	
2	<2020-06-03, cough, Clark>	0 (originally 1)	cross-symptom

Original database (with bounded contributions)

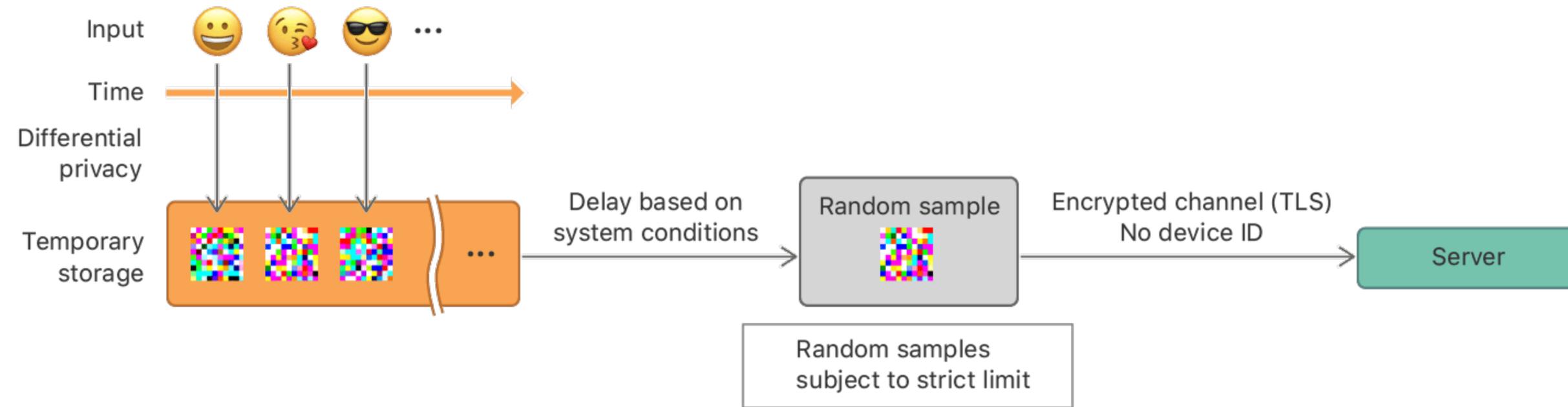
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# Local differential privacy



# Local differential privacy: Apple



# Count mean sketch with differential privacy (Apple)

Setup:

- Server samples a list of hash function  $H_1, H_2, \dots, H_k$
- Server initializes a matrix  $M$  of size  $m \times k$  to be all 0s

# Count mean sketch with differential privacy (Apple)

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To upload **val**, the client:

- Client samples a hash function  $H_i$  from list  $H_1, H_2, \dots, H_k$
- Client computes  $H_i(\mathbf{val}) = x$
- Client generates a one-hot vector of length  $m$  that is 0 everywhere and 1 at  $x$
- Client adds differential privacy to vector entries and sends the vector and  $i$  to the server
- The server adds the client's vector to row  $i$  of its matrix  $M$

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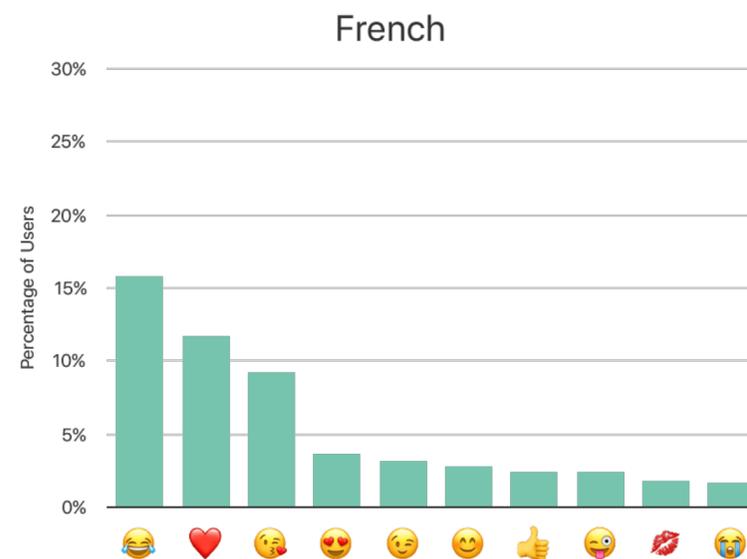
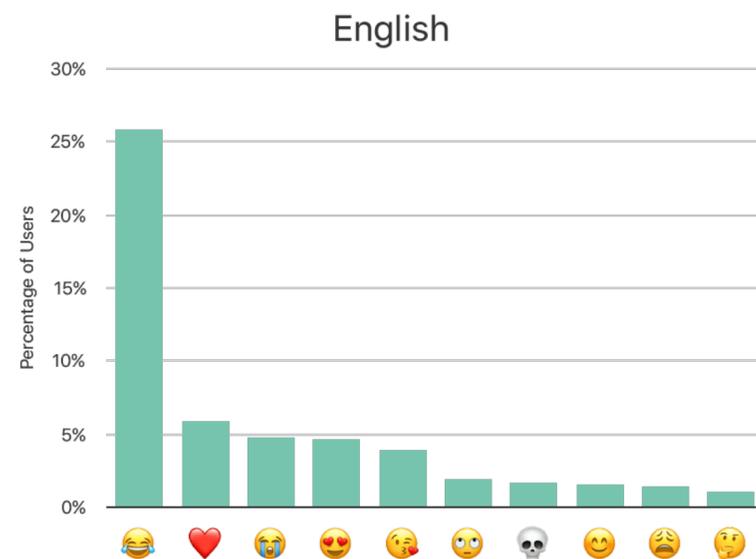
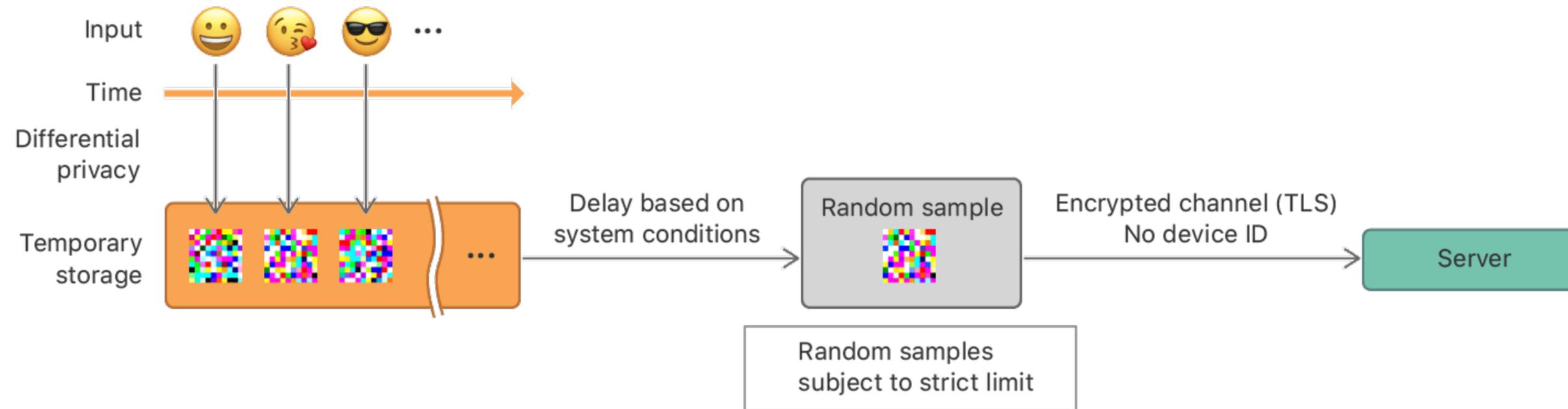
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To compute the mean of **val**:

- Server computes the average over  $M[i, H_i(\mathbf{val})]$  for  $i \in [k]$
- Note: requires the server to know the dictionary of domain elements

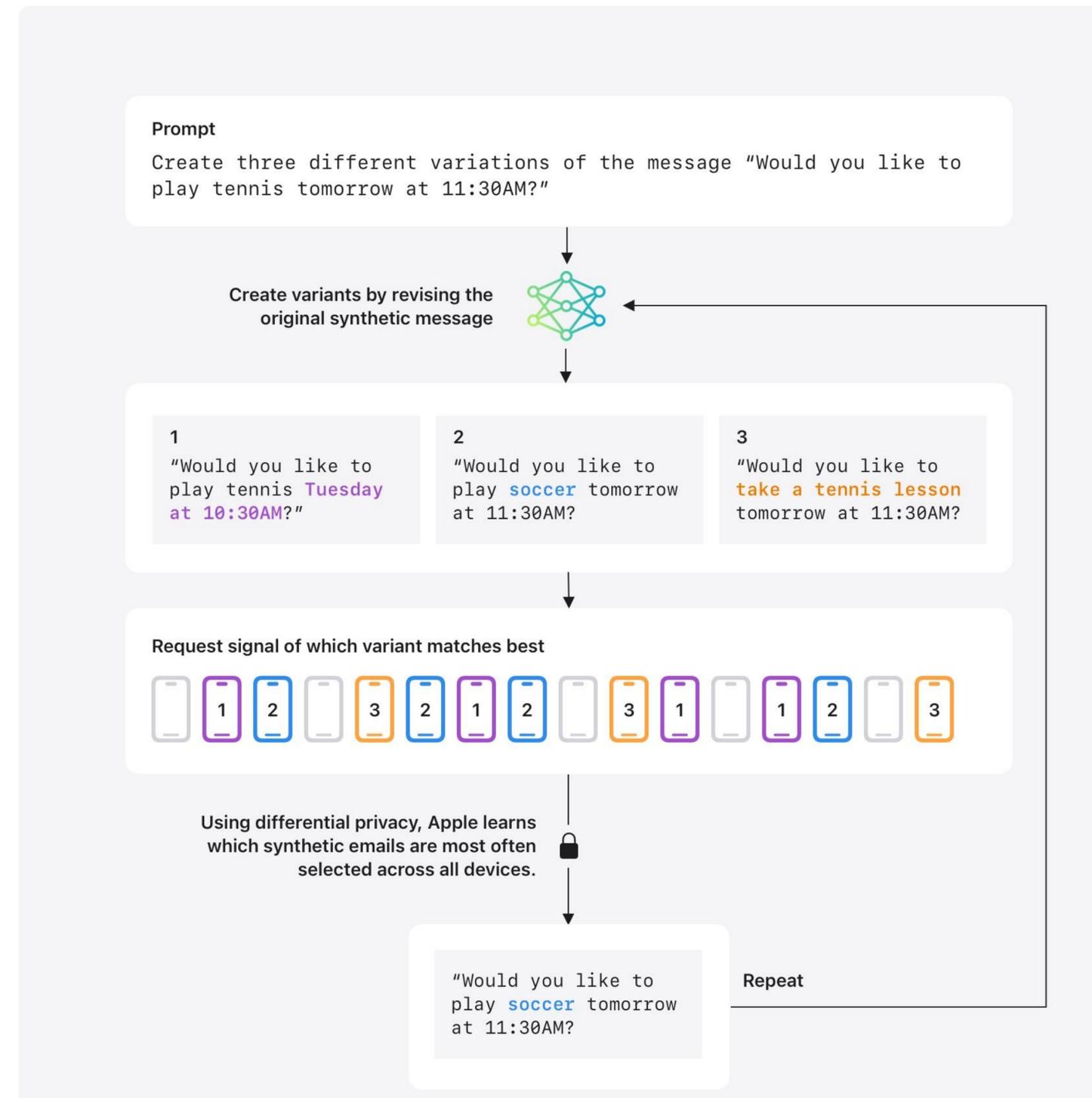
# Local differential privacy: Apple



(From 2017)

$$\epsilon = 4$$

# Local differential privacy: Apple



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# Logistics

Comments for project reports are in Gradescope

Final project presentations (12/2 and 12/4)

- 9 minute slots (hard cutoff time)
- Time for a few questions as the next group is setting up

Final project report (due 12/2)

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# References

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